



## Brief paper

Leader–follower containment control over directed random graphs<sup>☆</sup>Zhen Kan<sup>a</sup>, John M. Shea<sup>b</sup>, Warren E. Dixon<sup>a</sup><sup>a</sup> Department of Mechanical and Aerospace Engineering, University of Florida, Gainesville, USA<sup>b</sup> Department of Electrical and Computer Engineering, University of Florida, Gainesville, USA

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## ABSTRACT

The leader–follower consensus problem for multi-agent systems over directed random graphs is investigated. Motivated by the fact that inter-agent communication can be subject to random failure when agents perform tasks in a complex environment, a directed random graph is used to model the random loss of communication between agents, where the connection of the directed edge in the graph is assumed to be probabilistic and evolves according to a two-state Markov Model. In the leader–follower network, the leaders maintain a constant desired state and the followers update their states by communicating with local neighbors over the random communication network. Based on convex properties and a stochastic version of LaSalle's Invariance Principle, almost sure convergence of the followers' states to the convex hull spanned by the leaders' states is established for the leader–follower random network. A numerical simulation is provided to demonstrate the developed result.

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## 1. Introduction

Consensus problems that seek to agree upon certain quantities of interest have attracted significant research attention. A comprehensive review of consensus problems is provided in [Olfati-Saber, Fax, and Murray \(2007\)](#) and [Ren, Beard, and Atkins \(2007\)](#). To achieve consensus, agents are generally required to exchange information over a communication network as a means to coordinate their behaviors, such as achieving a common heading direction in flocking problems ([Jadbabaie, Lin, & Morse, 2003](#); [Tanner, Jadbabaie, & Pappas, 2007](#)), agreeing on the group average in distributed sensing ([Zhu & Martinez, 2010](#)), or achieving consensus in rendezvous and formation control problems ([Dimarogonas & Kyriakopoulos, 2007](#); [Kan, Navaravong, Shea, Pasiliao, & Dixon, 2015](#)), to name a few. In most of these applications, consistent information exchange between agents in either an undirected or directed manner is a common assumption to ensure full cooperation among team members. However, when agents operate in a complex environment, the inter-agent communication could be subject to

random failure due to either interference or unpredictable environmental disturbance. Since task completion relies on communication and interaction among agents, achieving consensus over such a stochastic communication network can be challenging.

Leader–follower containment control is a particular class of consensus problems, in which the networked multi-agent system consists of leader agents and follower agents. Generally, the leaders are a small subset of the agents, which are informed of the global task objectives, while the followers act under the influence of both neighboring agents and the leaders through local interactions. A main objective in leader–follower containment control is to drive all followers' states to a desired destination determined by the leaders' states. Hybrid control schemes are developed in [Ji, Ferrari-Trecate, Egerstedt and Buffa \(2008\)](#) to drive the dynamic follower agents into a convex polytope spanned by the stationary leader agents, where the local interaction among agents is modeled as an undirected graph. The work of [Ji et al. \(2008\)](#) is then extended to multiple stationary and dynamic leaders under a directed interaction graph in [Cao, Ren, and Egerstedt \(2012\)](#), [Li, Ren, and Xu \(2012\)](#) and [Meng, Ren, and You \(2010\)](#). Containment control for a leader–follower network under a switching graph is studied in [Lou and Hong \(2012\)](#) and [Notarstefano, Egerstedt, and Haque \(2011\)](#). In [Kan, Klotz, and Dixon \(2015\)](#), containment control is applied to a social network to regulate the emotional states of individuals to a desired end. For networked Lagrangian systems with parametric uncertainties, distributed containment control is developed in [Mei, Ren, and Ma \(2012\)](#). In the aforementioned works, a deterministic dynamic system is considered, where

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dynamic agents communicate and coordinate with other agents over an undirected or directed deterministic communication network. Mean-square containment control of a multi-agent system with communication noise is considered in Wang, Cheng, Hou, Tan, and Wang (2014). Since the results developed in Cao et al. (2012), Ji et al. (2008), Kan et al. (2015), Li et al. (2012), Lou and Hong (2012), Mei et al. (2012), Meng et al. (2010), Notarstefano et al. (2011), Wang et al. (2014) may not be applicable to stochastic communication networks where the existing communication links experience random loss, an extension of the classical containment control from the deterministic network to the stochastic network is desirable.

Building on graph theory and probability theory, several consensus results have been developed for random graphs. One of the earliest consensus results over an undirected random network is reported in Hatano and Mesbahi (2005), which proves that agreement can be achieved almost surely if the communication links between any pair of agents are activated independently with a common probability. The undirected random graph in Hatano and Mesbahi (2005) is extended to a general class of directed random graphs in Porfiri and Stilwell (2007) and Wu (2006). Necessary and sufficient conditions for consensus are developed in Tahbaz-Salehi and Jadbabaie (2008) for graphs that are generated by an ergodic and stationary random process. Mean-square-robust consensus over a network with communication noise and random packet loss is considered in the work of Zhang and Tian (2010) and Zhang and Tian (2012). Stochastic consensus for a multi-agent system with communication noise and Markovian switching topologies is investigated in Wang, Cheng, Ren, Hou, and Tan (2015). However, the convergence results reported in Hatano and Mesbahi (2005), Porfiri and Stilwell (2007), Tahbaz-Salehi and Jadbabaie (2008), Wang et al. (2015), Wu (2006), Zhang and Tian (2010), Zhang and Tian (2012) are only developed for leaderless networks without considering how the leaders can influence the followers to a desired end.

In this paper, the classical leader–follower containment control problem for deterministic systems is extended to a stochastic scenario. The leader–follower network is tasked to drive all followers into a prespecified destination area (i.e., the convex hull spanned by the leaders' states) under the influence of the leaders. Only the leaders are assumed to have the knowledge of the destination. To move toward the specified destination, the followers communicate and update their states with neighboring agents over a communication network. Since wireless communication is subject to random failure due to factors such as fading and packet loss, the inter-agent communication is modeled as a random graph, where each link evolves according to a two-state Markov Model to model the random loss of the existing communication link. In addition, the random communication network is assumed to be directed. Rather than assuming that all edges share a common edge probability and evolve independently with their previous edge connection states as in Hatano and Mesbahi (2005), different edges are allowed to have different transition probabilities in the current work that evolve according to a Markov Model, which can be used to model a large class of real-world networks to reflect the dependence of the current system states on their previous states. Moreover, compared to the works of Hatano and Mesbahi (2005), Porfiri and Stilwell (2007), Tahbaz-Salehi and Jadbabaie (2008), Wu (2006), a hierarchical network structure (i.e., leader–follower network) is considered where one-sided influence of leaders is used to affect the desired behaviors of the followers. Almost sure convergence of the followers' states to the convex hull spanned by the leaders' states over a random communication graph is then established via the convex properties in Boyd and Vandenberghe (2004) and a stochastic version of LaSalle's Invariance Theorem (Kushner, 1971).

## 2. Problem formulation

A multi-agent system consisting of  $n$  agents that communicate over wireless channels is considered. The wireless channels have intermittent connectivity, which cause the connections among the agents to vary with time. The communication graph is modeled as a temporal network, or time-varying graph,  $\mathcal{G}(t) = (\mathcal{V}, \mathcal{E}(t))$ . The vertices  $\mathcal{V}$  represent the agents, which do not vary with time. The edges  $\mathcal{E}(t)$  represent the connections among the agents and do vary with time. The flow of information is assumed to be asymmetric, so the edges in  $\mathcal{E}(t) \subset \mathcal{V} \times \mathcal{V}$  are directed. Specifically, the directed edge  $(v_j, v_i) \in \mathcal{E}$  indicates that node  $v_i$  can receive information from node  $v_j$ , but  $v_j$  may not necessarily receive information from  $v_i$ . In the directed edge  $(v_j, v_i)$ ,  $v_i$  and  $v_j$  are referred to as the child node and the parent node, respectively.

### 2.1. Directed random graph

Consider first the graph at one particular time, say  $t = t_0$ . Then, suppressing the time dependence,  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  is a directed graph. The graph  $\mathcal{G}$  is called a directed random graph if the set of edges  $\mathcal{E}$  is randomly determined. Let  $\tilde{\mathcal{E}} \subset \mathcal{V} \times \mathcal{V}$  be a set of potential directed edges connecting the nodes in  $\mathcal{V}$ . Each potential edge  $(v_j, v_i)$  is associated with a weight  $w_{ij} \in \mathbb{R}^+$ , which indicates how node  $v_i$  evaluates the information collected from  $v_j$ . We assume that the weight  $w_{ij}$  for each  $(v_j, v_i)$  is known initially and there are no self loops, so  $(v_i, v_i) \notin \tilde{\mathcal{E}}$ ,  $i = 1, 2, \dots, n$ . Associated with each potential edge  $(v_j, v_i) \in \tilde{\mathcal{E}}$ , let there be a Bernoulli random variable  $\delta_{ij}$ . An edge  $(v_j, v_i) \in \tilde{\mathcal{E}}$  will exist in  $\mathcal{E}$  if  $\delta_{ij} = 1$  and will not exist in  $\mathcal{E}$  if  $\delta_{ij} = 0$ . It is assumed that, for different edges, the  $\{\delta_{ij}\}$  are statistically independent.

Now, consider the temporal network,  $\mathcal{G}(t)$ , which consists of a time sequence of directed random graphs in which the edge set varies with  $t$ . In particular, each edge  $(i, j)$  evolves according to a two-state homogeneous Markov process  $\delta_{ij}(t)$  for  $i, j \in \{1, 2, \dots, n\}$  with stationary state transition probability  $p_{ij} \in (0, 1]$ , which indicates that, at the next time instant  $t'$ , the edge  $(i, j)$  will change its state to  $\delta_{ij}(t') = 1 - \delta_{ij}(t)$  with probability  $p_{ij}$  and will remain the previous state  $\delta_{ij}(t') = \delta_{ij}(t)$  with probability  $1 - p_{ij}$ .

**Assumption 1.** The random processes  $\{\delta_{ij}(t)\}$  do not change infinitely fast, and thus we can choose a sampling time  $\Delta_t$  such that with arbitrarily high probability,  $\delta_{ij}(t) = \delta_{ij}(t + t_0)$  if  $0 \leq t_0 < \Delta_t$  for all  $i, j \in \{1, 2, \dots, n\}$ .

Note that Assumption 1 will be true for any real system. For example, let  $T_0$  and  $T_1$  denote the expected dwell times in states 0 and 1 for the Markov process  $\delta_{ij}(t)$ , respectively. Then, the probability of staying in the same state during an observation period can be made arbitrarily large by selecting an appropriate  $\Delta_t$ . For example, the probability of remaining in state 0 during an interval of length  $\Delta_t$  is  $e^{-\Delta_t/T_0}$ .

We assume that the sequence of random graphs can be discretized in the following way. Let  $t_k = k\Delta_t$ ,  $k \in \mathbb{Z}^+$  be a time sequence, where  $\Delta_t \in \mathbb{R}^+$  is a sufficiently small sampling period during which we may assume the edge set is constant over each time interval  $[t_k, t_{k+1})$ . Let  $\mathcal{G}(k)$  denote the random graph  $\mathcal{G}(t)$  at  $t = t_k$ . Note that  $\mathcal{G}(k)$  is drawn from a finite sample space, which we denote by  $\tilde{\mathcal{G}} = \{\mathcal{G}_1, \dots, \mathcal{G}_M\}$ , and  $|\tilde{\mathcal{G}}| \leq 2^{|\tilde{\mathcal{E}}|}$ , which is determined by the power set of  $\tilde{\mathcal{E}}$ . In a directed graph, a *directed path* from node  $v_1$  to node  $v_k$  is a sequence of edges  $(v_1, v_2), (v_2, v_3), \dots, (v_i, v_k)$ . If a directed graph contains a *directed spanning tree*, every node has exactly one parent node except for one node, called the root, and the root has directed paths

to every other node. A directed graph is called *strongly connected* if there exists a directed path from every node to every other node in the graph. The weighted adjacency matrix  $A(k) = [a_{ij}(k)] \in \mathbb{R}^{n \times n}$  of the directed random graph  $\mathcal{G}(k)$  is

$$a_{ij}(k) = w_{ij} \delta_{ij}(k), \quad (1)$$

where  $w_{ij} \in \mathbb{R}^+$  is the edge weight. The Laplacian matrix  $L(k)$  is then defined as  $L(k) = D(k) - A(k)$ , where  $D(k) \triangleq [d_{ii}(k)] \in \mathbb{R}^{n \times n}$  is a diagonal matrix with the diagonal entry  $d_{ii}(k) = \sum_{j=1}^n a_{ij}(k)$ , and the off-diagonal entry  $d_{ij}(k) = 0$  for  $\forall i \neq j$ . Although  $A$  is a random matrix, the Laplacian matrix  $L$  is always a zero row sum matrix by its construction, which indicates that 0 is always an eigenvalue of  $L$  with the corresponding right eigenvector of  $\mathbf{1} = [1, \dots, 1]^T \in \mathbb{R}^n$ .

## 2.2. Leader–follower network over directed random graphs

Consider a multi-agent system composed of  $n$  agents that interact over a temporal network  $\mathcal{G}(t) = (\mathcal{V}, \mathcal{E}(t))$ , where, at each time, the graph  $\mathcal{G}(t)$  is a directed random graph, as described in Section 2.1. Suppose that the  $n$  agents in  $\mathcal{V}$  are partitioned into sets of  $\mathcal{V}_L$  with  $n_L \in \mathbb{Z}^+$  leaders and  $\mathcal{V}_F$  with  $n_F \in \mathbb{Z}^+$  followers. Let  $\mathbf{x}_i(k) \in \mathbb{R}^d$  denote the states (e.g., the Euclidean position) of agent  $v_i \in \mathcal{V}$  at the time instant  $t_k$ . Only the leaders are assumed to have the immutable and desired states that specifies the locations of the destination area where all agents are required to meet. The followers can only communicate with neighboring agents and update their states over the random communication graph  $\mathcal{G}(t)$ . Some directed graphs  $\mathcal{G}_i \in \bar{\mathcal{G}}$  can even be disconnected. Let  $\mathcal{G}_T = (\mathcal{V}, \bar{\mathcal{E}})$  be the graph that contains all possible edges for the graphs in  $\bar{\mathcal{G}}$ . To ensure the followers can be influenced by the leaders to the desired destination area over the directed random network, we introduce the following assumptions.

**Assumption 2.** It is assumed that  $\mathcal{G}_T \in \bar{\mathcal{G}}$  and  $\Pr\left(\bigcap_{j=k}^{k+n_F-1} \mathcal{G}(j) = \mathcal{G}_T\right) > 0$  for all  $k$ .

Note that Assumption 2 will be true in many real systems and practical models. For instance, it will be true for the Markov Model described after Assumption 1 if the Markov processes for the different edges are statistically independent. Such independent Markov Models can be used to model the effects of channel outages caused by multipath propagation. However, independence is not required and may not be present in some scenarios. For instance, if communication across the network is randomly blocked by a time-varying jammer, then the links may be completely dependent and yet Assumption 2 still holds.

**Assumption 3.** The graph  $\mathcal{G}_T \in \bar{\mathcal{G}}$  has a directed spanning tree, where, for each follower  $v_i \in \mathcal{V}_F$ , there exists at least one leader that has a directed path to the follower  $v_i$ .

Assumption 3 implies that the set of leaders act as the roots of the directed spanning tree in  $\mathcal{G}_T$ , which indicates the leaders have an influence directly or indirectly on all followers through a series of directed paths in the network. Different from the objective in Hatano and Mesbahi (2005) and Porfiri and Stilwell (2007), where every node has a probability to connect with every other node in the network, this work relaxes the constraint of having a strongly connected graph.

## 2.3. Objective

**Definition 1** (Boyd & Vandenberghe, 2004, Ch. 2). For a set of points  $z \triangleq \{z_1, \dots, z_n\}$ , the convex hull  $\text{Co}(z)$  is defined as the minimal

set containing all points in  $z$ , satisfying

$$\text{Co}(z) \triangleq \left\{ \sum_{i=1}^n \alpha_i z_i \mid z_i \in z, \alpha_i > 0, \sum_{i=1}^n \alpha_i = 1 \right\}.$$

**Definition 2** (Khalil, 2002, Ch. 4). A state  $z(k)$  approaches a set  $\mathcal{M}$  as  $t_k$  goes to infinity (i.e.,  $z(k) \rightarrow \mathcal{M}$  as  $t_k \rightarrow \infty$ ), if for each  $\varepsilon > 0$  there exists a  $T > 0$  such that  $\text{dist}(z(k), \mathcal{M}) < \varepsilon$  for  $t_k > T$ , where  $\text{dist}(z(k), \mathcal{M})$  denotes the distance from a point  $z(k)$  to a set  $\mathcal{M}$ . More precisely,

$$\text{dist}(z(k), \mathcal{M}) \triangleq \inf_{y \in \mathcal{M}} \|z(k) - y\|, \quad (2)$$

which is the smallest distance from  $z(k)$  to any point in  $\mathcal{M}$ .

To avoid notational confusion, let  $\mathbf{x}(k) \triangleq [x_1^T(k), \dots, x_n^T(k)]^T$  denote the stacked vector of all deterministic states  $x_i(k)$ ,  $v_i \in \mathcal{V}$ , and  $\mathbf{X}(k) \triangleq [\mathbf{X}_1^T(k), \dots, \mathbf{X}_n^T(k)]^T$  denote the random variables that represent the random states of  $v_i \in \mathcal{V}$  at time  $k$ . Let  $\mathbf{x}^F(k)$  and  $\mathbf{x}^L(k)$  denote the followers' states  $x_i(k)$ ,  $v_i \in \mathcal{V}_F$ , and the leaders' states  $x_i(k)$ ,  $v_i \in \mathcal{V}_L$ , respectively. The convex hull spanned by the states of leaders, and all states (i.e., both leaders and followers), are then represented as  $\text{Co}(\mathbf{x}^L(k))$  and  $\text{Co}(\mathbf{x}(k))$ , respectively. Since the leaders' states are static, the convex hull  $\text{Co}(\mathbf{x}^L(k))$  is constant, while the convex hull  $\text{Co}(\mathbf{x}(k))$  is time varying and depends on the states of the followers. The objective is to regulate the states of followers to a desired region, which is a convex hull spanned by all stationary leaders' states (i.e.,  $x_i(k) \rightarrow \text{Co}(\mathbf{x}^L(k))$ ), over a random communication network.

## 3. Consensus algorithms

Consider the random graph  $\mathcal{G}(k) = (\mathcal{V}, \mathcal{E}(k))$  at  $t_k$ . The follower updates its state  $\mathbf{X}_i(k)$ ,  $v_i \in \mathcal{V}_F$ , according to

$$\mathbf{X}_i(k+1) = \mathbf{X}_i(k) - \sum_{v_j \in \mathcal{N}_i(k)} K_g \Delta_t a_{ij}(k) (\mathbf{X}_i(k) - \mathbf{X}_j(k)), \quad (3)$$

where  $a_{ij}(k)$  is a random variable defined in (1),  $K_g \in \mathbb{R}^+$  is a control gain,  $\Delta_t$  is a small sampling time, and the time-varying set  $\mathcal{N}_i(k) \triangleq \{v_j \mid (v_j, v_i) \in \mathcal{E}(k)\}$  determines the set of the neighbors of  $v_i$  in  $\mathcal{G}(k)$ . Since the leaders are assumed to have desired constant states,

$$\mathbf{x}_i(k+1) = \mathbf{x}_i(k), \quad (4)$$

for  $\forall v_i \in \mathcal{V}_L$ .

Note that  $\mathbf{X}_i(k)$  is a random variable that evolves according to the stochastic system of (3). Let  $\{\mathbf{X}_i(k)\}$  be a Markov process. Since the systems in (3) and (4) along different dimensions are decoupled, for the simplicity of presentation,  $\mathbf{X}_i(k)$  will be treated as a scalar ( $d = 1$ ) in the subsequent analysis, and the extension of  $\mathbf{x}_i(k)$  to  $d$  dimensional states can be established by using the Kronecker product. The system of (3) and (4) can be rewritten in a compact form as

$$\begin{bmatrix} \mathbf{x}^L(k+1) \\ \mathbf{x}^F(k+1) \end{bmatrix} \triangleq \Phi(k) \begin{bmatrix} \mathbf{x}^L(k) \\ \mathbf{x}^F(k) \end{bmatrix} \quad (5)$$

where  $\Phi(k) \triangleq \begin{bmatrix} \mathcal{E}(k) \\ \pi(k) \end{bmatrix}$ ,  $\mathcal{E}(k) \triangleq [\mathbf{I}_{m \times m} \quad \mathbf{0}_{m \times (n-m)}] \in \mathbb{R}^{m \times n}$ ,  $\mathbf{I}_{m \times m} \in \mathbb{R}^{m \times m}$  denotes an identity matrix, and the entries of  $\pi(k) \in \mathbb{R}^{(n-m) \times n}$  are defined as

$$\pi_{il}(k) = \begin{cases} 1 - \sum_{v_j \in \mathcal{N}_i(k)} K_g \Delta_t a_{ij}(k) & i = l \\ \sum_{v_j \in \mathcal{N}_i(k)} K_g \Delta_t a_{ij}(k) & v_l \in \mathcal{N}_i(k), i \neq l \\ 0, & v_l \notin \mathcal{N}_i(k), i \neq l. \end{cases} \quad (6)$$



#### 4. Convergence analysis

In this section, almost sure convergence of the followers' states to the convex hull  $\text{Co}(x^L)$  spanned by the leaders' states is established for the agreement protocol in (5) over the directed random graph  $\mathcal{G}$ . To facilitate the subsequent convergence analysis, the definitions of almost sure convergence and supermartingale in a probabilistic setting are introduced.

**Definition 3** (Hatano & Mesbahi, 2005, Ch. 12). A random sequence  $\{\mathbf{Z}(k)\}$  in  $\mathbb{R}^n$  almost surely converges to a set  $M$ , if for every  $\epsilon > 0$

$$\lim_{k_0 \rightarrow \infty} \Pr \left\{ \sup_{k \geq k_0} \text{dist}(\mathbf{Z}(k), M) > \epsilon \right\} = 0, \quad (7)$$

where  $\text{dist}(\mathbf{Z}(k), M)$  is defined in (2) representing the distance of  $\mathbf{Z}(k)$  to the set  $M$ . Almost sure convergence is also called convergence with probability one (w.p.1).

**Definition 4** (Grimmett & Stirzaker, 2001). Let  $(\Omega, \mathcal{F}, \mathbb{P})$  be a measurable space. A filtration  $\mathcal{F}_0 \subseteq \mathcal{F}_1 \subseteq \mathcal{F}_2 \dots \subseteq \mathcal{F}_n$  is an increasing subsequence of sub- $\sigma$ -algebras of  $\mathcal{F}$ . A sequence of random variables  $\mathbf{Z}(k)$  is adapted to a filtration  $\mathcal{F}_k$  if  $\mathbf{Z}(k)$  is  $\mathcal{F}_k$ -measurable for all  $k$ . The pair  $(\mathbf{Z}, \mathcal{F})$  is called a supermartingale if, for all  $k \geq 0$ ,

$$\mathbb{E}[\mathbf{Z}(k)] < \infty \quad \text{and} \quad \mathbb{E}[\mathbf{Z}(k+1) | \mathcal{F}_k] \leq \mathbf{Z}(k), \quad (8)$$

where  $\mathbb{E}[\mathbf{Z}(k)]$  denotes the expected value of the random variable  $\mathbf{Z}(k)$ , and  $\mathbb{E}[\square | \diamond]$  denotes the conditional expectation of some  $\square$  under the condition of some  $\diamond$ .

The supermartingale sequence  $\{\mathbf{Z}(k)\}$  in (8) indicates that the current  $\mathbf{Z}(k)$  provides an upper bound for the conditional expectation  $\mathbb{E}[\mathbf{Z}(k+1) | \mathcal{F}_k]$  on the next time instant  $k+1$ , and  $\lim_{k \rightarrow \infty} \mathbf{Z}(k)$  exists and is finite w.p.1. In addition, if the sequence  $\{\mathbf{Z}(k)\}$  is a nonnegative supermartingale with  $\mathbb{E}[\mathbf{Z}(k)] < \infty$  in (8),  $\mathbf{Z}(k)$  converges w.p.1. to a limit (Kushner, 1971, Ch. 8).

##### 4.1. Almost sure convergence

In contrast to most results developed for the deterministic containment control (cf. Cao et al., 2012, Ji et al., 2008, Kan et al., 2015, Li et al., 2012, Lou & Hong, 2012, Meng et al., 2010, Notarstefano et al., 2011 and Mei et al., 2012) where network connectivity is the key assumption to ensure consensus (e.g., assuming a tree graph in directed or undirected graphs, or at least assuming a tree graph in the union of graphs for switching networks), the random setting in the current work does not make an explicit assumption on network connectivity at each time. The random graph  $\mathcal{G}(k)$  can be either connected or disconnected at each time  $t_k$ . However, below we show that Assumptions 2 and 3 ensure that the followers have sufficient access to the leaders' information so that the followers' states will almost surely converge to the convex hull spanned by the leaders' states.

Let  $V(x(k)) \in \mathbb{R}^+$  be the volume of the convex hull  $\text{Co}(x(k))$  spanned by all leaders' and followers' states at  $t_k$ . The strictly decreasing property of  $V(x(k))$  over a finite step is proven in Lemma 1 for the deterministic case of graph  $\mathcal{G}_T$ . For the stochastic case that every edge  $(v_j, v_i) \in \mathcal{E}$  in  $\mathcal{G}$  connects with a probability  $p_{ij} > 0$ , Lemma 2 indicates that  $V(\mathbf{X}(k))$  is nonincreasing at each time step. Based on Lemmas 1 and 2, almost sure convergence to the convex hull  $\text{Co}(x^L)$  for the stochastic system in (3) and (4) is then proven by using convex properties (Boyd & Vandenberghe, 2004) and the stochastic version of LaSalle's invariance principle in Kushner (1971, Ch. 8).

**Lemma 1.** Consider the particular directed graph  $\mathcal{G}_T$  where every potential edge in the random graph  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  is connected. Suppose

that Assumption 3 is satisfied, and the followers evolve according to (3), where the control gain  $K_g$  is selected sufficiently small such that  $\sum_{j \in \mathcal{N}_i} K_g \Delta_t a_{ij}(k) < 1$ . If there exists at least one node  $v_i \in \mathcal{V}_F$  with  $x_i(k) \notin \text{Co}(x^L)$ , the volume of the convex hull  $\text{Co}(x(k))$  will strictly decrease over  $n_f$  steps (i.e.,  $\text{Co}(x(k+n_f)) \subset \text{Co}(x(k))$  where  $n_f$  is the number of followers).

**Proof.** Since every potential edge exists in  $\mathcal{G}_T$ , the stochastic dynamics in (3) can be rewritten in a deterministic manner as

$$\begin{aligned} x_i(k+1) = & \left( 1 - \sum_{v_j \in \mathcal{N}_i} K_g \Delta_t a_{ij}(k) \right) x_i(k) \\ & + \sum_{v_j \in \mathcal{N}_i} K_g \Delta_t a_{ij}(k) x_j(k), \end{aligned} \quad (9)$$

for  $v_i \in \mathcal{V}_F$ . Here, the random set  $\mathcal{N}_i(k)$  in (3) is replaced by a constant set  $\mathcal{N}_i$  determined by  $\mathcal{G}_T$ . Since  $K_g$  is selected sufficiently small by assumption that  $\sum_{j \in \mathcal{N}_i} K_g \Delta_t a_{ij}(k) < 1$ ,  $x_i(k+1)$  is a convex combination of  $\{x_j(k)\}$ ,  $v_j \in v_i \cup \mathcal{N}_i$ . Note that the convex combination in (9) indicates that each vertex of  $\text{Co}(x(k))$  at  $k+1$  can either remain in the same state or evolve to shrink  $\text{Co}(x(k))$  by moving into the interior or along the boundary of  $\text{Co}(x(k))$ , resulting in  $\text{Co}(x(k+1)) \subseteq \text{Co}(x(k))$ . To show that the equality can be excluded in this set relation over  $n_f$  steps (i.e.,  $\text{Co}(x(k+n_f)) \subset \text{Co}(x(k))$ ), we will prove that at least one vertex of  $\text{Co}(x(k))$  evolves to shrink  $\text{Co}(x(k))$  over  $n_f$  steps. Since the existence of a node  $v_i \in \mathcal{V}_F$  with  $x_i(k) \notin \text{Co}(x^L)$  ensures that  $\text{Co}(x^L) \subset \text{Co}(x(k))$ , there always exists a vertex of  $\text{Co}(x(k))$  formed by followers only. Two cases for the component of such vertex are considered.

Case 1: Consider a vertex of  $\text{Co}(x(k))$  consisting of a single follower  $v_i \in \mathcal{V}_F$  such that  $x_i(k) \notin \text{Co}(x^L)$ . Since  $\mathcal{G}_T$  contains a directed path from  $\mathcal{V}_L$  to every node in  $\mathcal{V}_F$ , there exists at least one node  $v_j \in \mathcal{N}_i$  either in the interior or on the boundary of  $\text{Co}(x(k))$  with a different state  $x_j(k) \neq x_i(k)$ . According to (3), the vertex state  $x_i(k+1)$  will evolve either along the boundary or to the interior of  $\text{Co}(x(k))$  to shrink  $\text{Co}(x(k))$ , resulting in  $\text{Co}(x(k+1)) \subset \text{Co}(x(k))$ .

Case 2: Consider a vertex of  $\text{Co}(x(k))$  consisting of multiple followers  $\{v_i \in \mathcal{V}_F\}$  with the same states and each  $x_i(k) \notin \text{Co}(x^L)$ . Let  $S_v$  be the set of followers that forms such vertex. The worst case is that  $S_v$  contains all followers  $v_i$ ,  $i \in \{1, \dots, n_f\}$ , with  $x_1(k) = \dots = x_{n_f}(k)$ . According to (3), for those nodes  $v_i$  with  $\mathcal{N}_i \subset S_v$ ,  $x_i(k+1) = x_i(k)$ , due to  $x_i(k) = x_j(k)$ ,  $v_j \in \mathcal{N}_i$ . However,  $\mathcal{G}_T$  ensures a directed path from  $\mathcal{V}_L$  to at least one node  $v_j \in S_v$  with  $v_k \in \mathcal{N}_j$  and  $v_k \notin S_v$ . Following (3),  $x_j(k+1)$  moves out of the vertex either into the interior or along the boundary of  $\text{Co}(x(k))$ . Due to the connected  $\mathcal{G}_T$ , there always exists a node  $v_j$  in  $S_v$  with at least one neighbor  $v_k \in \mathcal{N}_j$  and  $v_k \notin S_v$ . Repeating the above process, the number of followers in  $S_v$  is strictly decreasing at each update, and Case 2 will reduce to Case 1 with the vertex containing a single follower for at most  $n_f$  steps, resulting in  $\text{Co}(x(k+n_f)) \subset \text{Co}(x(k))$  based on Case 1.  $\square$

**Lemma 2.** Let  $\mathcal{Q}_\lambda = \{x \in \mathbb{R}^n | V(x) \leq \lambda\}$  where  $\lambda > 0$ . Given that the sequence  $\{\mathbf{X}(k)\}$  evolves according to the dynamics in (3) and (4), if  $x(0)$  starts in  $\mathcal{Q}_\lambda$ ,  $\mathbf{X}(k) \in \mathcal{Q}_\lambda$  for all  $k \in \mathbb{Z}^+$  w.p.1., that is,

$$\Pr \left( \sup_{0 < k < \infty} V(\mathbf{X}(k)) > \lambda \right) = 0. \quad (10)$$

**Proof.** To show (10), it suffices to show that  $V(\mathbf{X}(k+1)) \leq V(\mathbf{X}(k))$ . Let  $\mathcal{N}_i(k)$  denote the set of all potential neighbors of  $v_i$  on  $\mathcal{G}_k$ . Consider a follower node  $v_i \in \mathcal{V}_F$  with a neighbor set  $\mathcal{N}_i(k)$

on  $\mathcal{G}_k$ , where  $\mathcal{N}_i(k) \subseteq \bar{\mathcal{N}}_i(k)$  indicates the set of nodes that are connected to  $v_i$  at  $t_k$ . Since each edge  $(v_j, v_i) \in \mathcal{E}_k$ ,  $v_j \in \mathcal{N}_i(k)$ , is connected with a probability  $p_{ij}$ , the node  $v_i$  may connect to either a subset of nodes in  $\bar{\mathcal{N}}_i$  or none of the nodes in  $\bar{\mathcal{N}}_i$ . If  $v_i$  connects to at least one node in  $\mathcal{N}_i(k)$ , according to Lemma 1,  $x_i(k)$  can either remain the same state or evolve to shrink the convex hull formed by itself and its neighbors  $v_j \in \mathcal{N}_i(k)$ , resulting in a nonincreasing  $V(x_k)$ . If  $v_i$  is isolated, which indicates that no edge exists between  $v_i$  and any other node in  $\mathcal{N}_i(k)$ ,  $x_i(k)$  will remain the same and not lead to an increasing volume of the convex hull  $\text{Co}(x(k))$ . Repeating this argument for every node in the graph  $\mathcal{G}_k$  indicates that  $V(\mathbf{X}(k+1)) \leq V(\mathbf{X}(k))$ .  $\square$

**Theorem 1.** Consider the random graph  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  that consists of stationary leaders and dynamic followers described by the stochastic system in (3) and (4). If Assumption 3 is satisfied, the followers  $v_i \in V_F$  will almost surely converge to the convex hull  $\text{Co}(x^L)$  that is spanned by the leaders states only.

**Proof.** The theorem is proven by using the stochastic version of LaSalle's invariance principle in Kushner (1971, Ch. 8) and convex properties in Boyd and Vandenberghe (2004, Ch. 2). Consider a time sequence  $T_m, m \in \mathbb{Z}^+$ , with  $T_{m+1} - T_m = n_f \Delta_t$  and  $T_0 = t_0$ . First, we show that there exists a positively invariant set  $\mathcal{Q}$  such that if  $x(0) \in \mathcal{Q}$ , then  $\mathbf{X}(T_m) \in \mathcal{Q}$  for all  $T_m \geq 0$ . Second, the volume of the convex hull  $\text{Co}(x(0))$  is decreasing to the invariant set  $E \subseteq \mathcal{Q}$  where the volume of  $\text{Co}(\mathbf{X}(T_m))$ ,  $\mathbf{X}(T_m) \in E$ , stays constant. It is then shown that  $M$  is the largest invariant set in  $E$ , where  $M$  is the set of points in the convex hull  $\text{Co}(x^L)$  formed by stationary leaders only.

Let  $\mathcal{Q} \triangleq \text{Co}(x(0))$ , where  $\text{Co}(x(0))$  is the convex hull formed by all initial states  $x(0)$ . Since Lemma 2 indicates that  $\mathbf{X}(k) \in \mathcal{Q}$  for all  $k \in \mathbb{Z}^+$  w.p.1. if  $x(0)$  starts in  $\mathcal{Q}$ ,  $\mathbf{X}(T_m) \in \mathcal{Q}$  also holds for all  $m \in \mathbb{Z}^+$  w.p.1., indicating that  $\mathcal{Q}$  is indeed a positively invariant set. To show that all the followers almost surely converge to the convex hull  $\text{Co}(x^L)$ , consider the auxiliary term  $\delta(x(T_m)) \in \mathbb{R}$  defined as

$$\delta(x(T_m)) \triangleq V(x(T_m)) - \mathbb{E}[V(\mathbf{X}(T_{m+1})) | \mathbf{X}(T_m) = x(T_m)]. \quad (11)$$

In (11), the state  $x(T_m)$  evolves according to (5) by following a sequence of  $n_f$  random graphs, where, by Assumption 1, each random graph stays constant for a period of  $\Delta_t$  over the interval  $[T_m, T_{m+1}] = n_f \Delta_t$ . To capture all possible sequences of the evolution from  $x(T_m)$  to  $x(T_{m+1})$ , let  $\bar{\mathcal{G}}' = \{\mathcal{G}'_1, \dots, \mathcal{G}'_{M_T}\}$  denote the finite set of all possible sequences over  $[T_m, T_{m+1}]$ , where  $M_T \in \mathbb{Z}^+$  denotes the cardinality of  $\bar{\mathcal{G}}'$  (i.e.,  $M_T = |\bar{\mathcal{G}}'|$ ) and each entry  $\mathcal{G}'_i \in \bar{\mathcal{G}}', i \in \{1, \dots, M_T\}$ , denotes a possible sequence that contains  $n_f$  graphs (i.e.,  $\mathcal{G}'_i \in \prod_{j=1}^{n_f} \bar{\mathcal{G}}$ ). Let  $p'_i \triangleq \Pr(\mathcal{G}'_i \in \bar{\mathcal{G}}')$  be the probability of the occurrence of the sequence  $\mathcal{G}'_i$  in  $\bar{\mathcal{G}}'$ .

Given the definition of the  $\bar{\mathcal{G}}'$  and the associated probability for each entry in  $\bar{\mathcal{G}}'$ , the conditional expectation in (11) is computed as

$$\mathbb{E}[V(\mathbf{X}(T_{m+1})) | \mathbf{X}(T_m) = x(T_m)] = \sum_{j=1}^{M_T} V(\Phi'_j x(T_m)) p'_j, \quad (12)$$

where  $\Phi'_j = \prod_{i=1}^{n_f} \Phi_i$  corresponds to the combined state transition matrix associated with the sequence  $\mathcal{G}'_j$  in  $\bar{\mathcal{G}}'$ , and each  $\Phi_i$  is the corresponding state transition matrix in (5) for the random graph at  $[t_{n_fm}, t_{n_fm} + i\Delta_t]$ . Let  $\mathcal{G}'_T$  be a path in  $\bar{\mathcal{G}}'$  that consists of  $\mathcal{G}_T$  only. By Assumption 2, the probability  $p'_T$  that  $\mathcal{G}_T$  occurs consecutively  $n_f$  times in  $\mathcal{G}'_T$  during  $[T_m, T_{m+1}]$  is strictly greater than zero,  $p'_T > 0$ . The strict decreasing property of  $V(\Phi'_j x(T_m)) p'_j$  is then established by Assumption 2 and Lemma 1, where  $\Phi'_T$  indicates the state transition matrix corresponding to  $\mathcal{G}'_T \in \bar{\mathcal{G}}'$ . For those entries  $\mathcal{G}'_j \in \bar{\mathcal{G}}'$  other than  $\mathcal{G}'_T$ , the convex properties ensure that each  $V(\Phi'_j x(T_m))$ ,  $j \in \{1, \dots, M_T\}$  and  $j \neq T$  in (12)

is a nonincreasing function as shown in Lemma 2. Hence, when considering all possible graphs in  $\bar{\mathcal{G}}$ ,

$$\mathbb{E}[V(\mathbf{X}(T_{m+1})) | \mathbf{X}(T_m) = x(T_m)] < V(x(T_m)),$$

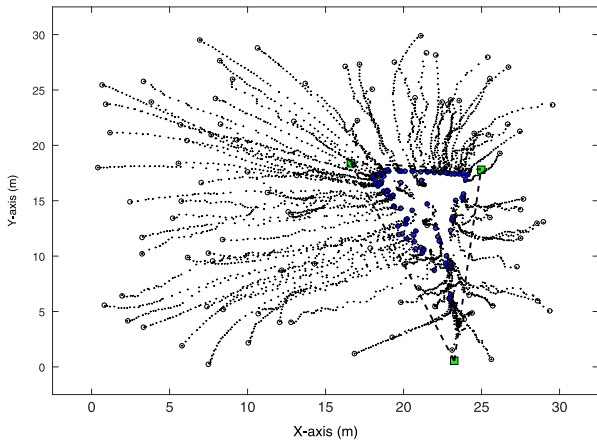
which indicates that the sequence of  $\{V(\mathbf{X}(T_m))\}$  is a supermartingale. Since  $\{V(\mathbf{X}(T_m))\}$  is a nonnegative supermartingale and bounded in  $\mathcal{Q}$ , based on Definition 4 and Theorem 1 in Kushner (1971, Ch. 8),  $V(\mathbf{X}(T_m))$  will decrease to a limit, which indicates that the convex hull  $\text{Co}(\mathbf{X}(k))$  will shrink to an invariant set as  $m \rightarrow \infty$ .

Let  $E$  be the invariant set of all points in  $\mathcal{Q}$ , where the volume of  $\text{Co}(\mathbf{X}(T_m))$ ,  $\mathbf{X}(T_m) \in E$ , stays constant. Based on LaSalle's Theorem, the largest invariant set  $M$  must be established, where  $M$  is the set of points in the convex hull  $\text{Co}(x^L)$  formed by stationary leaders only. A proof by contradiction is used to show that  $M$  (i.e.,  $\text{Co}(x^L)$ ) is the largest invariant set. Let  $M' \supset M$  be a larger invariant set in  $E$ . Suppose that there is a follower whose state  $x_i \notin M$ , and  $x_i$  is on the boundary of  $M'$ , while the other followers  $x_j \in M$ ,  $v_j \in \mathcal{V}_F - \{v_i\}$ . Since  $M' \subset E$ , the volume of the set  $M'$  stays constant. The only way for the volume of  $M'$  to stay constant is that  $x_i(t)$  stays constant for all  $t \geq 0$ . However, for this to happen, the follower  $v_i$  must be isolated from the group for all  $t \geq 0$ . This isolation is a contradiction with Assumption 2. Hence,  $M$  is the largest invariant set. A stochastic version of LaSalle's invariance principle in Kushner (1971, Ch. 8) can now be invoked to ensure almost sure convergence of the followers' states to the largest invariant set  $M$  (i.e., the convex hull  $\text{Co}(x^L)$ ).  $\square$

**Remark 1.** Containment control for a multi-agent system has received significant focus (cf. Cao et al., 2012, Ji et al., 2008, Kan et al., 2015, Li et al., 2012, Lou & Hong, 2012, Mei et al., 2012, Meng et al., 2010, Notarstefano et al., 2011 and Wang et al., 2014). However, few results investigate collision avoidance among agents in containment control over either deterministic or stochastic communication networks. In our previous work (Kan et al., 2015), containment control is applied to a deterministic social network to regulate the emotional states of individuals to a desired end. To maintain the social bond within a desired threshold (i.e.,  $\|q_i(t) - q_j(t)\| \leq \delta$ , where  $q_i$  and  $q_j$  denote the time-varying social states of individual  $i$  and  $j$ , respectively, and  $\delta \in \mathbb{R}^+$  denotes the threshold), a navigation function based framework from Kan, Dani, Shea, and Dixon (2012) is developed to ensure existing social influence between individuals during network evolution. Since ensuring collision avoidance among agents (i.e.,  $\|x_i(t) - x_j(t)\| \geq \xi$ , where  $x_i$  and  $x_j$  denote the time-varying positions of agent  $i$  and  $j$ , respectively, and  $\xi \in \mathbb{R}^+$  denotes a safe inter-agent distance) is analogous to maintaining existing social influence between individuals, the approach developed in Kan et al. (2015) could be modified to handle collision avoidance among agents in containment control. In another recent work (Cheng, Kan, Rosenfeld, Parikh, & Dixon, 2014), a navigation function based distributed controller is developed for a multi-agent system to perform formation control with ensured collision avoidance over an intermittent sensing network where the agents suffer random loss of inter-agent sensing and communication. Based on the stochastic analysis framework developed in the present work, the approach developed in Cheng et al. (2014) could also be extended for collision avoidance in containment control over stochastic communication networks.

## 5. Simulation

A graph of 100 nodes in a two-dimensional coordinate space is considered, where the states of nodes are their locations in the two-dimensional coordinate space. The followers and the leaders form a directed graph, represented by the dots and squares in

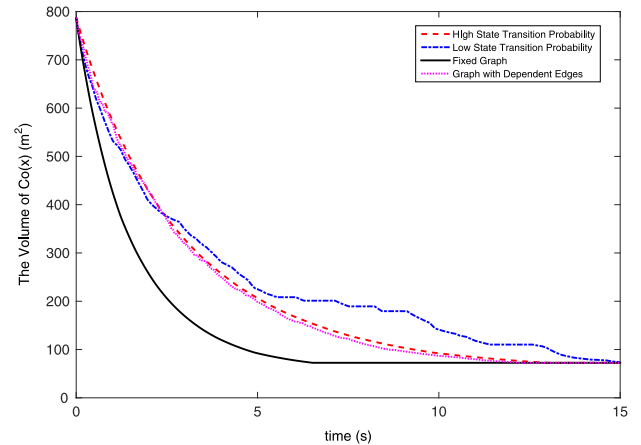


**Fig. 1.** Node trajectories plot. The followers' states converge to the convex hull formed by stationary leaders, where the leaders' states are denoted as squares, and the followers' initial and final states are denoted by circles and dots, respectively.

**Fig. 1.** respectively. To show random loss of the communication links, a two-state Markov Model with independent edges is applied to represent the evolution of the edge states. Particularly, each edge  $(i, j)$  will change its current state (i.e., from connection to disconnection or from disconnection to connection) with probability  $p_{ij}$  and will remain in the current state (i.e., connection or disconnection) with probability  $1 - p_{ij}$ . Consider the case that  $p_{ij} \in (0, 1)$  is randomly generated for each edge in the graph. **Fig. 1** indicates that the followers converge to the locations inside the convex hull spanned by the stationary leaders on the given random graph. We also compared the performance of the same graph with different connection models and the results are shown in **Fig. 2**. In the comparison, all graphs start from the same initial graph and differ in how the edge connections are evolved according to the two-state Markov Model. The solid line in **Fig. 2** indicates the evolution of  $\text{Co}(x)$  for a fixed connection model where all potential edges are connected all the time. The dashed and dot-dashed lines in **Fig. 2** represent the evolution of  $\text{Co}(x)$  for the Markov Model with independent edges that each edge  $(i, j)$  fails or succeeds with high state transition probability  $p_{ij} \in [0.95, 1)$  and low state transition probability  $p_{ij} \in (0, 0.05]$ , respectively. The dot line shows a Markov Model with dependent edges that all fail or succeed at the same time, where the stationary state probabilities are generated randomly from a uniform  $(0, 1)$  distribution. As expected, the performance of the fixed graph outperforms other connection models, since constant information exchange is available all the time. However, such a fully connected graph appears rarely in a random network that evolves according to the two-state Markov Model (e.g., appearing only 4 times for a simulation of 15,000 times). Note that the algorithm developed in the current work can achieve the same containment result without requiring the graph to be regularly fully connected, as shown in **Fig. 2**. The graph with dependent edges and the graph with high state transition probability behave similar to each other, converging to the convex hull  $\text{Co}(x^L)$  at about the same speed. The graph with low state transition probability takes the longest time to converge, since the graph may run into a disconnected graph and stay there for a relatively long time due to its low probability in state transition, as shown by the non-strictly decreasing dot-dashed line.

## 6. Conclusion

Leader–follower containment control on random graphs is examined. The underlying random network is assumed to evolve according to a two-state Markov Model, and the followers are proven



**Fig. 2.** The evolution of the volume of the convex hull  $\text{Co}(x)$  for the graph with different connection models.

to converge to the convex hull spanned by the stationary leaders w.p.1. The simulation results demonstrate the convergence of all followers to the convex hull  $\text{Co}(x^L)$  over random networks. Although convergence on random graphs is established, convergence speed is not investigated in the current work. Since the convergence speed of consensus highly relies on the topology of the underlying communication network, additional work will examine the topology design to increase the convergence speed for the developed leader–follower containment control on random graphs.

## Acknowledgment

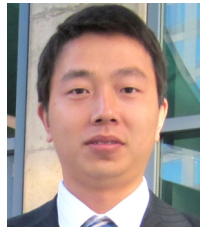
Any opinions, findings and conclusions or recommendations expressed in this material are those of the authors and do not necessarily reflect the views of the sponsoring agency.

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